

# E-ACSL and future Frama-C plug-in

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**HI-LITE**

Simplifying the use of formal methods



# Part I

## E-ACSL

long ra  
for 0 <=  
C1); if (m  
tmp2 =  
of the

tmp2[j] = 0; if (i <= (Nbr - 1)) else if (tmp1[j] >= 0) { i <= (Nbr - 1); tmp2[j] = (i <= (Nbr - 1) ? 1 : 0); else tmp2[j] = tmp1[j]; } /\* Then the second part looks like the first one: \*/  
tmp1[i][k] = 0; k <= 5; k++) tmp1[i][k] += m2[i][k] \* tmp2[k]; } /\* The [i,j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
i <= 3; tmp1[i][i] >= 1; \*/ Final rounding: tmp2[i][i] is now represented on 9 bits: \*if (tmp1[i][i] < -256) m2[i][i] = -256; else if (tmp1[i][i] > 255) m2[i][i] = 255; else m2[i][i] = tmp1[i][i];



## Executable Ansi/ISO C Specification Language

What should be?

- ▶ executable subset of ACSL
- ▶ preserve ACSL semantics as much as possible
- ▶ compatible with ALFA as much as possible

Which goals?

- ▶ runtime assertion checking
- ▶ usable by dynamic analyses tools
- ▶ usable by static verification tools like Frama-C plug-ins
- ▶ verification of mixed ADA/C programs



- ▶ last version: 1.5-3
- ▶ [Hi-Lite deliverable 3.4.1](#)
- ▶ based on ACSL v1.5
- ▶ detailed syntax
- ▶ mainly point out differences with ACSL



- ▶ similar to ACSL (e.g. mathematical integer arithmetic)
- ▶ 3-valued logic with **undefinedness**
  - ▶  $1/0 == 1/0$  is undefined
  - ▶  $f(*p)$  is undefined if  $p$  is invalid
  - ▶ tools must ensure that undefined terms are never evaluated
- ▶ **lazy operators**  $\&\&, ||, \_? \_ : \_, ==>$ 
  - ▶  $\backslash \text{false} \&\& 1/0 == 1/0$  is invalid
  - ▶  $1/0 == 1/0 \&\& \backslash \text{false}$  is undefined
  - ▶ different but consistent semantics compared to ACSL
  - ▶ a valid (resp. invalid) E-ACSL predicate is valid (resp. invalid) in ACSL



- ▶ 3 different kinds of quantifications (only 1 in ACSL)
- ▶ unguarded quantification *à la* ACSL only allowed for boolean and char
- ▶ guarded integer quantification

```
\forall typ x1, ..., xn;
  a1 <= x1 <= b1 ... && an <= xn <= bn
  ==> p
```

- ▶ guarded **iterator** quantification
  - ▶ from which element does the iteration begin?
  - ▶ how to access to the next elements?
  - ▶ which guards must be true to continue to iterate?
  - ▶ the only syntax extension from ACSL



```
struct btree {
    int val;
    struct btree *left, *right;
};

/*@ iterator access(_, struct btree *t):
    @   nexts t→left, t→right;
    @   guards \valid(t→left),
           \valid(t→right); */

/*@ predicate is_even(struct btree *t) =
    @   \forall struct btree *tt;
           access(tt, t) ==> tt→val % 2 == 0; */
```



- ▶ lose their inductive nature
- ▶ a loop invariant  $I$  is equivalent to
  - ▶ put an **assertion  $I$**  just before entering the loop
  - ▶ put the same **assertion** at the very end of the loop body

long n  
for (i = 0; i < n; i++)  
c[i] = 0;  
tmp2 = 0;  
for (k = 0; k < n; k++)  
for (j = 0; j < n; j++)  
tmp2[i][j] = 0;

tmp2[i][j] = 0; // i < (n-1) also if tmp1[i][j] >= 0 // i < (n-1) tmp2[i][j] = (i < (n-1) ? 0 : tmp2[i][j]) + tmp1[i][j]; // Then the second part looks like the first one. // tmp1[i][j] = 0; k = 0; k < n; k++) tmp1[i][j] += m2[i][k] \* tmp2[k][j]; // The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1 \*MC1 // i = 0; tmp1[i][j] >= 0; // Final rounding: tmp2[i][j] is now represented on 9 bits. \*if (tmp1[i][j] < -256) tmp2[i][j] = -256; else if (tmp1[i][j] > 255) tmp2[i][j] = 255; else tmp2[i][j] = tmp1[i][j];



## Present:

- ▶ recursive logic definitions
- ▶ specification modules

## Absent:

- ▶ lemmas and axiomatic (not computable)
- ▶ inductive predicate (not computable in general)
- ▶ polymorphism and higher order (still experimental in ACSL)
- ▶ concrete logic type (still experimental in ACSL)
- ▶ memory footprint (still experimental in ACSL)



# Part II

## Future Frama-C Plug-in

long ra  
for 0 ->  
C1) if (m  
tmp2 =  
re of the

tmp2[j] = 0; for (k = 0; k < (Nb1 - 1); k++) tmp2[j] += m2[k][j] \* tmp1[k]; /\* The [j] coefficient of the matrix product MC2\*TMP1, that is \*MC2\*(MC1\*M1) = MC2\*(M1\*MC1) = 1. tmp1[0][i] >= -1.\*/ Final rounding: tmp2[0][i] is now represented on 9 bits. \*if (tmp1[0][i] < -256) m2[0][i] = -256; else if (tmp1[0][i] > 255) m2[0][i] = 255; else m2[0][i] = tmp1[0][i];



- ▶ new Frama-C plug-in called 'E-ACSL'
- ▶ takes a C program annotated with ACSL as input
- ▶ **checks** that annotations are part of E-ACSL
- ▶ roughly **converts annotations**

```
int div(int x, int y) {
  /*@ assert y != 0; */
  return x / y;
}
```

into C code

```
int div(int x, int y) {
  /*@ assert y != 0; */
  if (y == 0) e_acsl_fail();
  return x / y;
}
```



- ▶ E-ACSL integers are mathematical integers
- ▶ heavy translation via **GMP** (could be optimized)

```

/*@ assert -3 == x; */ ;
// declare temp variables
mpz_t tmp1, tmp2, tmp3, int tmp 4;
mpz_init_set_si(tmp1,3); // init tmp1 with 3
mpz_init(tmp2); // init tmp2
mpz_neg(tmp2, tmp1); // tmp2 = -tmp1 = -3
mpz_init_set_si(tmp3, x); // init tmp3 with x
// really check the assertion by comparing -3 to x
tmp4 = mpz_cmp(tmp2, tmp3);
if (tmp4 != 0) e_acsl_fail("(-3 == x)");
// deallocate temp variables
mpz_clear(tmp1); mpz_clear(tmp2); mpz_clear(tmp3);
  
```



in Hi-Lite, only expect to handle a big-enough subset of E-ACSL

## unsupported features:

- ▶ floats and reals
- ▶ constructs for memory management like `\valid` or `\at`

## partially supported features:

- ▶ memory accesses: will not check that they are valid

```

long n;
for (i = 0; i < n; i++)
  tmp2 =
  of the

```

```

tmp2[i] = 1 + (n-i) * i; // also #tmp1[i] >= 1 + (n-i) * i; // #tmp2[i] = tmp1[i]; // Then the second part looks like the first one
tmp1[i] = 0; k = 5; k--> tmp1[i] = mc2[i][k] * tmp2[k]; // The [i][k] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1 *MC1
i = 1; tmp1[0] >= 1; // Final rounding: tmp2[0] is now represented on 9 bits: #if (tmp1[0] < -255) m2[0] = -255; else #if (tmp1[0] > 255) m2[0] = 255; else #endif

```



- ▶ yet preliminary development
- ▶ first version planed **to be released in September/October**
- ▶ not possible to implement the whole stuff
- ▶ **need your feedback** to know what E-ACSL features you are going to use
- ▶ will implement required features before others



long n...  
 for 0...  
 C1) if...  
 tmp2...  
 of the...  
 tmp2[0] = 1 << (n-1) + abs(tmp1[0]) <> 1 << (n-1) + abs(tmp1[0]) <> (n-1) + abs(tmp1[0]) <> tmp1[0]; /\* Then the second part takes the first one...  
 tmp1[0] = 0; k = 5; k--> tmp1[0] += m2[0][k] \* tmp2[k]; /\* The [j] coefficient of the matrix product MC2\*TMP2, that is \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1 + MC1...  
 l = 1; tmp1[0] >>= 1; /\* Final rounding: tmp2[0] is now represented on 9 bits. If (tmp1[0] < -255) tmp1[0] = -255; else if (tmp1[0] > 255) tmp1[0] = 255; else tmp1[0] =...

Any questions?

long ra  
for 0 ->  
C1); if (m  
tmp2 =  
se of the

tmp2[j][i] = 0; if (i < (Nbr - 1)) else if (tmp1[j][i] >= 0) tmp2[j][i] = (1 << (Nbr - 1)) - 1; else tmp2[j][i] = tmp1[j][i]; /\* Then the second part takes like the first one: \*/  
tmp1[i][k] = 0; k = 0; k <= 5; k++) tmp1[i][k] += m2[j][k] \* tmp2[k][i]; /\* The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(M1\*M1) = MC2\*M1\*M1  
i = 1; tmp1[i][i] >= -1; /\* Final rounding: tmp2[i][i] is now represented on 9 bits: \*if (tmp1[i][i] < -256) m2[i][i] = -256; else if (tmp1[i][i] > 255) m2[i][i] = 255; else m2[i][i] = tmp1[i][i];

