

# E-ACSL Frama-C plug-in

Julien Signoles

CEA LIST

Software Safety Labs

Hi-Lite Meeting

November 2011, the 29<sup>th</sup>



**HI-LITE**

Simplifying the use of formal methods



## Executable Ansi/ISO C Specification Language

<http://www.open-do.org/wp-content/uploads/2011/05/e-acsl.pdf>

### What is it?

- ▶ executable subset of ACSL
- ▶ preserve ACSL semantics as much as possible
- ▶ compatible with ALFA as much as possible

### Which goals?

- ▶ runtime assertion checking
- ▶ usable by dynamic analyses tools
- ▶ usable by static verification tools like Frama-C plug-ins
- ▶ verification of mixed ADA/C programs



- ▶ takes an **annotated C program** as input
- ▶ checks that each annotation belongs to E-ACSL
- ▶ **returns a new C program**
- ▶ equivalent to the input
- ▶ **each annotation is converted into new C statements**
- ▶ including (at least) one guard
- ▶ which fails at runtime if the annotation is wrong

long n;  
for (i = 0; i < n; i++)  
C[i] = 0;  
tmp2 =  
of the

tmp2[i][j] = i \* (n-1) + j; tmp1[i][j] = i \* (n-1) + j; tmp2[i][j] = tmp1[i][j];  
tmp1[i][j] = 0; k = 5; k = tmp1[i][j] \* m2[i][k] + tmp2[i][k];  
Final rounding: tmp2[i][j] is now represented on 9 bits: if (tmp1[i][j] < -255) m2[i][j] = -255; else if (tmp1[i][j] > 255) m2[i][j] = 255; else m2[i][j] = tmp1[i][j];



▶ input:

```
int div(int x, int y) {
    /*@ assert y != 0; */
    return x / y;
}
```

▶ output:

```
int div(int x, int y) {
    /*@ assert y != 0; */
    if (y == 0) e_acsl_fail();
    return x / y;
}
```

▶ a correct translation is much more complicated



- ▶ use **GMP integers** when required
- ▶ **keep the annotation** for documentation and further uses
- ▶ usually **one block** of statements **by annotation** (not always possible, e.g. \at)
- ▶ inserted at the **right code location**
- ▶ declares **temporary variables**
  - ▶ memoization to reduce memory usage
  - ▶ at function/global level when required
- ▶ **allocates and deallocates** them when required
- ▶ contains a **guard** if (! guard) e\_acsl\_fail(msg);
- ▶ may contain additional guards to **prevent execution of undefined values** (or at least a warning right now)



```
/*@ assert y != 0; */ z = x / y;
```

1. push a new environment `env` to translate the annotation
2. translate term `y` of type `int` to the `int` expression `y`
3. coerce `y` to an integer
  - 3.1 generate a fresh `mpz_t` variable `e_acsl_1` corresponding to `y`
  - 3.2 add its declaration to `env`
  - 3.3 add its initialisation to `env`
    - 3.3.1 as the type of `y` is signed and smaller than `long`, generate `mpz_init_set_si(e_acsl_1, y);`
  - 3.4 add its deallocation to `env`
    - 3.4.1 generate `mpz_clear(e_acsl_1);`
  - 3.5 translate `y` to `e_acsl_1`
4. translate term `0` of type `integer` to a fresh `mpz_t` variable `e_acsl_2`



5. as its operands are **integers**, **translate** `!=` by using `mpz_cmp`
  - 5.1 generate a **fresh int variable** `e_acsl_3`
  - 5.2 add its declaration to `env`
  - 5.3 add its initialisation to `env`
    - 5.3.1 generate `e_acsl_3 = mpz_cmp(e_acsl_1,e_acsl_2);`
  - 5.4 no deallocation of `e_acsl_3` required
  - 5.5 translate `y != 0` to `e_acsl_3 != 0`
6. **add the guard** checking the assertion to `env`
  - 6.1 `e_acsl_3 != 0` already gets type `int`: right!
  - 6.2 add the statement `if (! (e_acsl_3 != 0)) then e_acsl_fail("y != 0");` to `env`
7. **extend** `/*@ assert y != 0; */ z = x / y;` with a new **block** computed from `env` and `z = x / y;`
8. `pop env`



- ▶ option `-e-acsl` to run the plug-in
- ▶ resulting code put in a new Frama-C project "e-acsl"
- ▶ new code **linkable against GMP**
- ▶ new code **analysable by other analysers**
- ▶ use standard Frama-C options on these projects
- ▶ option `-e-acsl-project` to set the resulting project name

Demo!





### implemented

- ▶ C types
- ▶ integer
- ▶ boolean
- ▶ implicit coercions

### not yet implemented

- ▶ real

```

(long n)
for (i = 0; i < n; i++)
  tmp2 =
  ...

```

```

tmp2[i] = (i < (n-1)) ? tmp1[i] : (i < (n-1) ? tmp2[i] : tmp1[i]);
// Then the second pass takes like the first pass:
tmp1[0] = 0; k = 5; k--; tmp1[k] += m2[k] * tmp2[k]; // The [k] coefficient of the matrix product M2*TMP2, that is: *M2*(TMP1) = M2*(M1*M1) = M2*M1*M1
i = i - tmp1[0] >= 1; // Final rounding: tmp2[0] is now represented on 9 bits: *if (tmp1[0] < -256) m2[0] = -256; else if (tmp1[0] > 255) m2[0] = 255; else tmp2[0] =

```



### implemented

- ▶ integer constants
- ▶ C left values
- ▶ arithmetic operators
- ▶ casts
- ▶ address &
- ▶ sizeof
- ▶ alignof
- ▶ `\null` (as `(void *)0`)
- ▶ `\at` (extra restriction)
- ▶ `\result`

### not yet implemented

- ▶ `\true` and `\false`
- ▶ bitwise operators
- ▶ boolean operators
- ▶ conditional
- ▶ let binding
- ▶ typedef
- ▶ t-sets
- ▶ `\base_addr`, `\offset` and `\block_length`



### implemented

- ▶ `\true` and `\false`
- ▶ relations (`==`, `<=`, ...)
- ▶ lazy conjunction `&&`
- ▶ lazy disjunction `||`
- ▶ lazy implication `==>`
- ▶ negation `!`

### not yet implemented

- ▶ equivalence `<==>`
- ▶ exclusive or `^^`
- ▶ conditionals
- ▶ let bindings
- ▶ quantifications
- ▶ `\at`
- ▶ `\valid et al.`
- ▶ `\initialized`



### implemented

- ▶ assertions
- ▶ function contracts
- ▶ statement contracts

### not yet implemented

- ▶ behavior-specific annotations
- ▶ loop annotations
- ▶ global annotations

```

long n;
for (i = 0; i < n; i++)
  tmp2[i] = 0;

```

```

tmp2[0] = 1; // (n-1) else if (tmp1[j]) >= 1; // (n-1) - 1; else tmp2[j] = tmp1[j]; // Then the second part looks like the first one.
tmp1[0] = 0; k = 5; k--> tmp1[0][j] = mc2[0][k] * tmp2[k][j]; // The [j] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*MC1
i = 1; tmp1[0][i] >= 1; // Final rounding: tmp2[0][i] is now represented on 9 bits. *if (tmp1[0][i] < -256) m2[0][i] = -256; else if (tmp1[0][i] > 255) m2[0][i] = 255; else tmp2[0][i] =

```



### implemented

- ▶ assumes
- ▶ requires
- ▶ ensures

### not yet implemented

- ▶ assigns
- ▶ decreases
- ▶ abrupt termination
- ▶ complete behaviors
- ▶ disjoint behaviors

long n  
for 0 <=  
C1; if (n  
tmp2 =  
of the

tmp2[0] = 1; for (int i = 1; i < n; i++) tmp2[i] = 1; for (int i = 1; i < n; i++) tmp2[i] = 1; Then the second part takes like the first one.  
tmp1[0] = 0; for (int k = 1; k <= n; k++) tmp1[k] = mc2[0][k] \* tmp2[k]; // The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
i = 1; tmp1[0] >= 1; // Final rounding, tmp2[0] is now represented on 9 bits. if (tmp1[0] < -255) tmp1[0] = -255; else if (tmp1[0] > 255) tmp1[0] = 255; else tmp1[0] = tmp1[0];



- ▶ release of first prototype planned for January 2012
  - ▶ based on **Frama-C Nitrogen-20111001**
  - ▶ implement some missing useful features
    - ▶ quantifiers over integers
    - ▶ what else?
  - ▶ plug-in packaging and documentation
  - ▶ stronger testing
- ▶ new release of **E-ACSL reference manual** at the same time
- ▶ use case during 2012
- ▶ implement other missing useful features during 2012
- ▶ better handling of **E-ACSL undefined terms**
  - ▶ will require Frama-C Oxygen
- ▶ improve customizability on need
- ▶ **internship proposal: executable C memory model**



Any questions?

long ra  
for 0 ->  
C1); if (m  
tmp2 =  
se of the

tmp2[j][i] = 0; if (i < (Nbr - 1)) else if (tmp1[j][i] >= 0) tmp2[j][i] = (1 << (Nbr - 1)) - 1; else tmp2[j][i] = tmp1[j][i]; /\* Then the second part takes like the first one: \*/  
tmp1[i][k] = 0; k = 0; k <= 5; k++) tmp1[i][k] += mC2[i][k] \* tmp2[k][j]; /\* The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(M1)\*M1 = MC2\*M1\*M1  
i = 1; tmp1[i][i] >= -1; /\* Final rounding: tmp2[i][i] is now represented on 9 bits. \*if (tmp1[i][i] < -256) m2[i][i] = -256; else if (tmp1[i][i] > 255) m2[i][i] = 255; else m2[i][i] = tmp1[i][i];

